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Charliecloud: Unprivileged Containers for User-Defined Software Stacks Title:

in HPC

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Charliecloud

Unprivileged Containers for User-Defined Software Stacks in HPC



Tim Randles Reid Priedhorsky

November 15, 2017



The next 25 minutes of your life

- 1. Why user-defined software stacks will end your suffering
- 2. But only if you use containers
- 3. Use Charliecloud and all your wildest dreams will come true

Some people need different software

Default software stacks are good at specific things.

– in the case of HPC, it's MPI-based simulation codes

What if your thing is different?

- non-MPI simulations
- data analytics and machine learning
- epic build process

Admins will install software for you.

- BUT only if there's enough demand
- unusual needs go unmet
- are you crackpot or innovative?

Solution: User-defined software stacks

BYOS (bring your own software)

- Let users install software of their choice
- ... up to and including a complete Linux distribution
- ... and run this image on compute resources they don't own.



Why User-Defined Software Stacks (UDSS)?

Advantages

- software dependencies: numerous, unusual, older, newer, internet ...
- portability of environments: e.g., across dev/test/small/large ...
- consistent environments: validated, standardized, archival ...
- usability

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Disadvantages (possibly)

- missing functionality: HSN, accelerators, file systems
- performance: many opportunities for overhead



- 1. Standard, reproducible workflow
- 2. Work well on existing resources
- 3. Be very simple

Design goals

1. Standard, reproducible workflow

- in contrast with "tinker 'til it's ready, then freeze"
- standard ⇒ reduce training/devel costs, increase skill portability
- reproducible ⇒ creation of images is easier & more robust

2. Work well on existing resources

- HPC centers are very good at what they do
- let's not re-implement and re-optimize
 resource management: solved (Slurm, Moab, Torque, PBS, etc.)
 file systems: solved (Lustre, Panasas, GPFS)
 high-speed interconnect: solved (InfiniBand, OPA)

3. Be very simple

- save costs: development, debugging, security, usability
- UNIX philosophy: "make each program do one thing well"

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options	definition	kernel	core libraries	app libraries	pros	cons

UDSS shares with host							
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compile it yourself	download all your dependencies and compile them	yes	yes	mixed	always available; in principle can do anything	not 1995 anymore; in practice too hard	

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Full-featured

- image building
- image management storage, caching, tagging, signing
- orchestration
- storage management
- runtime setupe.g., default command/script, inetd-alike
- stateful containers
- supervisor daemon(s)

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systemd-nspawn [???] NsJail [???]

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1450ali [! !]

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- 1. code size
- 2. support burden
- 3. privileged & trusted operations

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Conclusion: Lightweight implementations are a better choice for HPC centers

- most important cloud-like flexibility
- don't compromise existing tools & strengths of HPC centers

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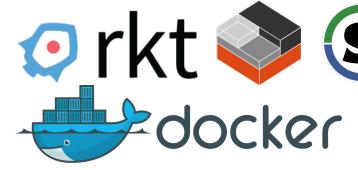
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But ... some of those other features are important

1. Linux namespaces

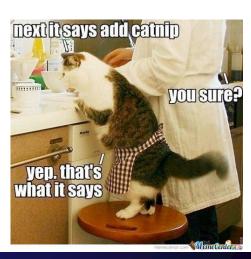
- mount: filesystem tree and mounts
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- UTS: host name & domain name
- network: all other network stuff
- IPC: System V and POSIX
- user: UID/GID/capabilities

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2. cgroups

- limit resource consumption per process
- 3. prctl (PR_SET_NO_NEW_PRIVS)
 - prevent execve(2) from increasing privileges
- 4. seccomp(2)
 - filter system calls
- 5. SELinux, AppArmor, etc.
 - various features that change what a process may do



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Charliecloud's hybrid approach

1. Image building & sharing goes in a sandbox

- safe place for users to be root: user workstation or virtual machine
- use Docker for image building or anything else that can produce a filesystem tree debootstrap(8), yum --installroot, etc.
- wrap Docker for image management ch-docker2tar

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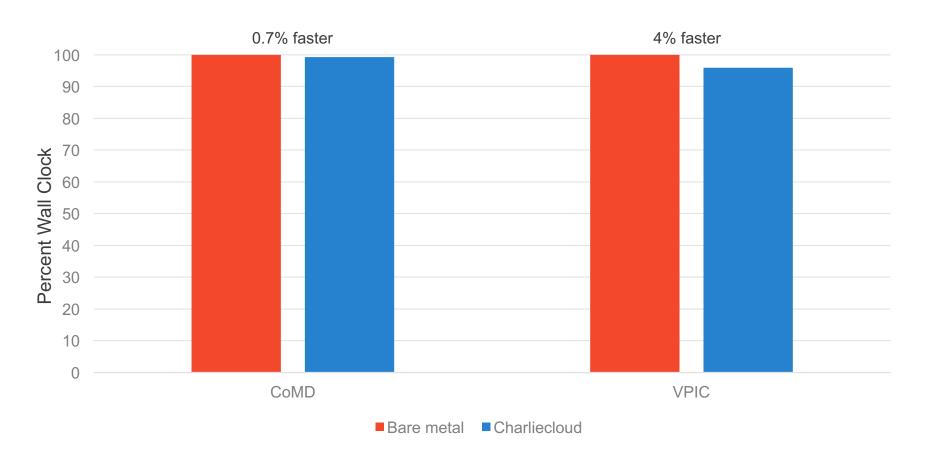
2. Run images with our own unprivileged runtime

- mount & user namespaces only
 - requires new-ish kernel
 - most distros have the right kernel (Fedora in 2015, Ubuntu Xenial in 2016)
 - Cray UP04 has it
 - RHEL/CentOS 7 can install via ElRepo (or enable on kernel command line in 7.4)
- it's a user program!!!
- admins don't need to do anything

Basic workflow

step	wl	here	privileged?	
Step	sandbox	production	privilegea:	
1. Build Docker/etc. image	✓		maybe	
2. Dump image to tarball	✓		maybe	
3. Copy tarball to where you want to run	✓	✓	no	
4. Unpack tarball		✓	no	
5. Configure your stuff (sometimes)		✓	no	
6. Run your commands in container		✓	no	

Performance e.g.: CoMD and VPIC (32 nodes)



Charliecloud vs. the design goals

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Charliecloud status

1. Available now on some LANL clusters

- passes tests on Crays
- Woodchuck (IC) now, Fog (ASC) very soon

2. Installable now on any Linux box

- newer kernel needed (roughly 4.4+)
- including cloud instances

3. Instructions for pre-installed VirtualBox image

- no root needed
- Mac, Windows, Linux, Solaris

4. Packages available on openSUSE Build Service (community)

- CentOS 7, Debian 9.0, Xubuntu 16.04 & 17.10
- 5. PR for HTCondor integration (community)

Charliecloud resources

;login: article (USENIX magazine)

- "Linux containers for fun and profit in HPC"
- https://www.usenix.org/publications/login/fall2017/priedhorsky

Supercomputing 2017

- "Charliecloud: Unprivileged containers for UDSS in HPC"

Documentation

- https://hpc.github.io/charliecloud
- includes detailed tutorials

Source code

https://github.com/hpc/charliecloud

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